

MyDUL 测试流程和结果

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测试步骤

普通表测试

在本节的测试中, 我们创建两个 Oracle 数据库用户, 一个是 andy, 另一个是 jake, 我们在用户 andy 下创建 4 张表, 通过系统的数据字典视图插入数据。然后通过 mydul 程序把用户 andy 中的表导出来, 然后通过 myimp 程序把数据装载到用户 jake 下, 然后比较两个用

户下数据的差异，测试 mydul 程序导出数据的准确性。

创建表空间

使用缺省块大小为 8K 的表空间。

```
create tablespace ts_dul datafile 'ts_dul_01.dbf' size 100M autoextend on;
```

创建用户

使用上面创建的表空间做为用户的缺省表空间。

```
create user andy identified by andy default tablespace ts_dul;
```

为用户授权，使得用户能使用资源和查询任意数据字典

```
grant connect, resource to andy;
```

```
grant select any dictionary to andy;
```

```
create user jake identified by jake default tablespace ts_dul;
```

```
grant connect, resource to jake;
```

创建测试表

连接到用户 andy，通过系统视图创建普通表。

```
conn andy/andy
```

```
CREATE TABLE MYOBJECT AS SELECT * FROM DBA_OBJECTS;
```

```
CREATE TABLE MYPRIV AS SELECT * FROM DBA_SYS_PRIVS;
```

```
CREATE TABLE MYSPACE AS SELECT * FROM DBA_TABLESPACES;
```

```
CREATE TABLE MYTABLE AS SELECT * FROM DBA_TABLES;
```

连接到 Oracle 系统用户，执行 checkpoint 把数据写入到数据库文件中

```
conn / as sysdba
```

```
alter system checkpoint;
```

创建后,表中已经有数据了,通过 mydul 导出用户 andy 下的数据存储成文件,通过 myimp 导入到另一个用户 jake 中,通过 minus 比较两个用户下的数据是否一致。

导出用户数据

编辑 mydul 的配置文件,把系统表空间数据文件和刚创建的 ts_dul 表空间数据文件加进去。

```
[road@CentOS etc]$ cat config.ini  
  
dictdir=/home/road/mydul/dict  
  
datadir=/home/road/mydul/data  
  
datafiles=/home/road/mydul/etc/dbfiles.list
```

```
[road@CentOS etc]$ cat dbfiles.list  
  
/oracle11/oradata/orallg/system01.dbf  
  
/oracle11/oradata/orallg/users01.dbf  
  
/oracle11/product/11.2.0/dbs/road_ts01.dbf  
  
/oracle11/product/11.2.0/dbs/ts_dul_01.dbf
```

在执行目录 bin 下运行

```
[road@CentOS bin]$ ./mydul config=../etc/config.ini
```

```
MYDUL> export dict  
MYDUL> load dict  
MYDUL> unload user andy
```

导入用户数据

转到数据存放目录,看到导出文件名为 ANDY.dat 的数据文件,通过 myimp 装载到用户 jake 下。

```
[road@CentOS bin]$ cd ../data
```

```
[road@CentOS data]$ ls
ANDY.dat
[road@CentOS data]$ ../bin/myimp jake/jake file=ANDY.dat touser=jake
```

比较数据一致性

连接到系统用户

```
conn / as sysdba
```

```
SQL> select * from andy.myobject minus select * from jake.myobject;
no rows selected
```

```
SQL> select * from andy.mypriv minus select * from jake.mypriv;
no rows selected
```

```
SQL> select * from andy.myspace minus select * from jake.myspace;
no rows selected
```

```
SQL> select * from andy.mytable minus select * from jake.mytable;
no rows selected
```

如果上面的两个用户中表的数据有不一致的地方，就会显示不同的数据行出来。

结论

原始数据与 mydul 导出的数据一致，测试通过。

cluster 表测试

本节的测试我们创建两张 cluster 成员表，为了方便测试，并且能有做够数据，我们创建的表还是从数据字典视图中获取数据。用户仍然使用 andy 和 jake。

创建 cluster

连接到用户 andy。

```
conn andy/andy
```

```
create cluster c_object_id (object_id number);
create index i_object_id on cluster c_object_id;
```

创建成员表

连接到用户 andy。

```
conn andy/andy
```

```
create table clu_object
(
  owner          varchar2(30),
  object_name    varchar2(30),
  object_id      number,
  data_object_id number,
  object_type    varchar2(19),
  created        date,
  status         varchar2(7)
)
cluster c_object_id (object_id);
```

```
create table clu_table
(
  object_id      number,
  owner          varchar2(30),
  table_name     varchar2(30),
  tablespace_name varchar2(30),
  cluster_name   varchar2(30),
  status         varchar2(8),
  pct_free       number,
  pct_used       number,
  ini_trans      number,
  max_trans      number,
  logging        varchar2(3),
  degree         varchar2(10),
  partitioned    varchar2(3)
)
cluster c_object_id (object_id);
```

在成员表中插入数据

连接到用户 andy。

```
conn andy/andy
```

```
insert into clu_object
select
  owner, object_name, object_id, data_object_id,
```

```
    object_type, created, status
from dba_objects;

commit;

insert into clu_table
select
    o.object_id, t.owner, t.table_name, t.tablespace_name,
    t.cluster_name, t.status, t.pct_free, t.pct_used,
    t.ini_trans, t.max_trans, t.logging, t.degree, t.partitioned
from dba_objects o, dba_tables t
where o.owner = t.owner and o.object_name = t.table_name;

commit;
```

checkpoint 把数据刷新到磁盘中

```
SQL> conn / as sysdba
Connected.
SQL> alter system checkpoint;
System altered.
```

在第一个成员表 clu_object 中插入数据时，操作过程及其缓慢，这是由于 Oracle 要为 cluster 的块预留空间，造成磁盘空间消耗非常快，要不断分配 extent。所以数据量大的时候用 cluster 表是不明智的。上面的操作，如果要减小数据量，可以把 SYS 用户的对象剔除掉。

导出表数据

```
[road@CentOS bin]$ ./mydul config=../etc/config.ini
```

```
MYDUL> export dict
MYDUL> load dict
MYDUL> unload table andy.clu_object
MYDUL> unload table andy.clu_table
```

导入表数据

```
[road@CentOS bin]$ cd ../data
[road@CentOS data]$ ls
```

```
ANDY_CLU_OBJECT.dat ANDY_CLU_TABLE.dat
[road@CentOS data]$ ../bin/myimp jake/jake file=ANDY_CLU_OBJECT.dat touser=jake

[road@CentOS data]$ ../bin/myimp jake/jake file=ANDY_CLU_TABLE.dat touser=jake
```

注意 cluster 表在目标端创建后，成为普通表。

比较数据一致性

连接到系统用户

```
conn / as sysdba
```

```
SQL> select
  2   object_id, owner, object_name, data_object_id,
  3   object_type, created, status
  4 from andy.clu_object
  5 minus
  6 select
  7   object_id, owner, object_name, data_object_id,
  8   object_type, created, status
  9 from jake.clu_object;
no rows selected
```

```
SQL> select * from andy.clu_table minus select * from jake.clu_table;
no rows selected
```

第一个表由于建表时的字段顺序与源表不一致，所以要把查询的字段按一致的顺序排列才能进行比较。

多键 cluster

同上面的测试步骤，只是把 cluster 键设为 owner 和 object_name。

连接到用户 andy。

```
conn andy/andy
```

创建 cluster

```
create cluster c_owner_name (owner varchar2(30), object_name varchar2(30));
create index i_owner_name on cluster c_owner_name;
```


创建成员表

```
create table c_object
(
  owner          varchar2(30),
  object_name    varchar2(30),
  object_id      number,
  data_object_id number,
  object_type    varchar2(19),
  created        date,
  status         varchar2(7)
)
cluster c_owner_name (owner, object_name);
```

```
create table c_table
(
  owner          varchar2(30),
  table_name     varchar2(30),
  tablespace_name varchar2(30),
  cluster_name   varchar2(30),
  status         varchar2(8),
  pct_free       number,
  pct_used       number,
  ini_trans      number,
  max_trans      number,
  logging        varchar2(3),
  degree         varchar2(10),
  partitioned    varchar2(3)
)
cluster c_owner_name (owner, table_name);
```

插入数据

```
insert into c_object
select
  owner, object_name, object_id, data_object_id,
  object_type, created, status
from dba_objects where owner <> 'SYS';
commit;
```

```
insert into c_table
select
  owner, table_name, tablespace_name,
  cluster_name, status, pct_free, pct_used,
  ini_trans, max_trans, logging, degree, partitioned
```

```
from dba_tables
where owner <> 'SYS';
commit;
```

```
checkpoint 把数据刷新到磁盘中
SQL> conn / as sysdba
Connected.
SQL> alter system checkpoint;
System altered.
```

导出成员表数据

```
[road@CentOS bin]$ ./mydul config=../etc/config.ini
```

```
MYDUL> export dict
MYDUL> load dict
MYDUL> unload table andy.c_object
MYDUL> unload table andy.c_table
```

导入成员表数据

```
[road@CentOS bin]$ cd ../data
[road@CentOS data]$ ls
ANDY_C_OBJECT.dat  ANDY_C_TABLE.dat
[road@CentOS data]$ ../bin/myimp jake/jake file=ANDY_C_OBJECT.dat touser=jake

[road@CentOS data]$ ../bin/myimp jake/jake file=ANDY_C_TABLE.dat touser=jake
```

注意 cluster 表在目标端创建后，成为普通表。

比较数据一致性

连接到系统用户

```
conn / as sysdba
```

```
SQL> conn / as sysdba
Connected.
```

```
SQL> select * from andy.c_object minus select * from jake.c_object;
```

```
no rows selected
```

```
SQL> select * from andy.c_table minus select * from jake.c_table;
```

```
no rows selected
```

用户 andy 和 jake 下的表 c_object 和 c_table 数据一致。

结论

cluster 成员表的原始数据与 mydul 导出的数据一致，测试通过。

删除 cluster

有两种方法，一种是按创建的相反顺序，把对象依次删除。

```
drop table clu_object purge;
drop table clu_table purge;
drop index i_object_id;
drop cluster c_object_id;
```

第二种方法是删除 cluster 连带一起删除成员表。

```
drop cluster c_object_id including tables;
```

行依赖表的测试

本节测试还在用户 andy 下进行。仍然从视图中获取数据。

创建表

连接到用户 andy，通过系统视图创建行依赖表。

```
conn andy/andy
```

```
create table rd_object rowdependencies as select * from dba_objects where rownum<1;
```

插入数据

```
declare
  v_obj dba_objects%ROWTYPE;
  cursor c_obj is select * from dba_objects;
begin
  open c_obj;
  loop
    fetch c_obj into v_obj;
    exit when c_obj%NOTFOUND;
    insert into rd_object values (v_obj.owner, v_obj.object_name,
```

```
v_obj.subobject_name, v_obj.object_id, v_obj.data_object_id,
v_obj.object_type, v_obj.created, v_obj.last_ddl_time,
v_obj.timestamp, v_obj.status, v_obj.temporary,
v_obj.generated, v_obj.secondary, v_obj.namespace,
v_obj.edition_name
);
commit;
end loop;
close c_obj;
end;
/
```

行依赖表中的每条数据的 ora_rowscn 是数据提交时的 SCN，为了使每行的 SCN 不一样，我们每插入一条数据，进行一次提交。

查询出 20 行看一下，结果如下：

```
SQL> select ora_rowscn, object_id from rd_object where rownum < 21;
```

ORA_ROWSCN	OBJECT_ID
4093190	85
4093192	86
4093194	87
4093196	88
4093198	89
4093200	90
4093202	91
4093203	92
4093205	93
4093207	94
4093209	95

ORA_ROWSCN	OBJECT_ID
4093211	96
4093213	97
4093215	98
4093217	99
4093218	100
4093219	101
4093220	102
4093221	103
4093223	104

20 rows selected.

checkpoint 把数据刷新到磁盘中

```
SQL> conn / as sysdba
```

Connected.

```
SQL> alter system checkpoint;
```

System altered.

导出表数据

```
[road@CentOS bin]$ ./mydul config=../etc/config.ini
```

```
MYDUL> export dict
```

```
MYDUL> load dict
```

```
MYDUL> unload table andy.rd_object
```

导入表数据

```
[road@CentOS bin]$ cd ../data
```

```
[road@CentOS data]$ ls
```

ANDY_RD_OBJECT.dat

```
[road@CentOS data]$ ../bin/myimp jake/jake file=ANDY_RD_OBJECT.dat touser=jake
```

比较数据一致性

```
SQL> select * from andy.rd_object minus select * from jake.rd_object;
```

no rows selected

结论

原始数据与 mydul 导出的数据一致，测试通过。

在数据导入建表的过程中，只建立了表中简单的列关系，表的其他属性都忽略了，如果需要创建表的其他属性，可以手动建表，在导入数据时忽略建表错误即可。

从下面看到，用户 jake 下的 rd_object 表的 rowdependencies 属性没有了。

```
SQL> select dependencies from dba_tables
```

```
where owner='ANDY' and table_name='RD_OBJECT';
```

```
DEPENDEN
```

```
-----
```

```
ENABLED
```

```
SQL> select dependencies from dba_tables  
where owner='JAKE' and table_name='RD_OBJECT';
```

```
DEPENDEN
```

```
-----
```

```
DISABLED
```

unused 列测试

本节测试中仍然使用 andy 用户，从 dba_objects 视图中获取数据，数据分两部分，一部分是正常表时的数据，然后 unuse 一列，把剩余数据插入。

创建表

```
SQL> select count(*) from dba_objects;
```

```
COUNT(*)
```

```
-----
```

```
71998
```

```
SQL> select count(*) from dba_objects where object_id < 20000;
```

```
COUNT(*)
```

```
-----
```

```
19514
```

先创建表，把 object_id 小于 20000 的数据导入进去。

```
create table un_object as select * from dba_objects where object_id < 20000;
```

unuse column

```
SQL> desc un_object;
```

```
Name                               Null?    Type
```

OWNER	VARCHAR2 (30)
OBJECT_NAME	VARCHAR2 (128)
SUBOBJECT_NAME	VARCHAR2 (30)
OBJECT_ID	NUMBER
DATA_OBJECT_ID	NUMBER
OBJECT_TYPE	VARCHAR2 (19)
CREATED	DATE
LAST_DDL_TIME	DATE
TIMESTAMP	VARCHAR2 (19)
STATUS	VARCHAR2 (7)
TEMPORARY	VARCHAR2 (1)
GENERATED	VARCHAR2 (1)
SECONDARY	VARCHAR2 (1)
NAMESPACE	NUMBER
EDITION_NAME	VARCHAR2 (30)

```
SQL> alter table un_object set unused column timestamp;
```

Table altered.

```
SQL> desc un_object;
```

Name	Null?	Type
OWNER		VARCHAR2 (30)
OBJECT_NAME		VARCHAR2 (128)
SUBOBJECT_NAME		VARCHAR2 (30)
OBJECT_ID		NUMBER
DATA_OBJECT_ID		NUMBER
OBJECT_TYPE		VARCHAR2 (19)
CREATED		DATE
LAST_DDL_TIME		DATE
STATUS		VARCHAR2 (7)
TEMPORARY		VARCHAR2 (1)
GENERATED		VARCHAR2 (1)
SECONDARY		VARCHAR2 (1)
NAMESPACE		NUMBER
EDITION_NAME		VARCHAR2 (30)

插入剩余数据

```
insert into un_object  
select owner, object_name, subobject_name, object_id,
```

```
data_object_id, object_type, created, last_ddl_time,  
status, temporary, generated, secondary, namespace, edition_name  
from dba_objects where object_id >= 20000;
```

```
commit;
```

checkpoint 把数据刷新到磁盘中

```
SQL> conn / as sysdba
```

```
Connected.
```

```
SQL> alter system checkpoint;
```

```
System altered.
```

导出表数据

```
[road@CentOS bin]$ ./mydul config=../etc/config.ini
```

```
MYDUL> export dict
```

```
MYDUL> load dict
```

```
MYDUL> unload table andy.un_object
```

导入表数据

```
[road@CentOS bin]$ cd ../data
```

```
[road@CentOS data]$ ls
```

```
ANDY_UN_OBJECT.dat
```

```
[road@CentOS data]$ ../bin/myimp jake/jake file=ANDY_UN_OBJECT.dat touser=jake
```

比较数据一致性

```
SQL> conn / as sysdba
```

```
Connected.
```

```
SQL> select * from andy.un_object minus select * from jake.un_object;
```

```
no rows selected
```

结论

原始数据与 mydul 导出的数据一致，测试通过。

dropped 列测试

本节的测试步骤与 unused 列测试相似，只是在准备数据时把列 drop 掉。仍然使用 andy 用户。

数据准备

```
SQL> conn andy/andy
```

```
Connected.
```

```
SQL> create table dr_object as select * from dba_objects where object_id < 20000;
```

```
Table created.
```

```
SQL> desc dr_object
```

Name	Null?	Type
OWNER		VARCHAR2(30)
OBJECT_NAME		VARCHAR2(128)
SUBOBJECT_NAME		VARCHAR2(30)
OBJECT_ID		NUMBER
DATA_OBJECT_ID		NUMBER
OBJECT_TYPE		VARCHAR2(19)
CREATED		DATE
LAST_DDL_TIME		DATE
TIMESTAMP		VARCHAR2(19)
STATUS		VARCHAR2(7)
TEMPORARY		VARCHAR2(1)
GENERATED		VARCHAR2(1)
SECONDARY		VARCHAR2(1)
NAMESPACE		NUMBER
EDITION_NAME		VARCHAR2(30)

```
SQL> alter table dr_object drop column temporary;
```

```
Table altered.
```

```
SQL> desc dr_object;
```

Name	Null?	Type
OWNER		VARCHAR2(30)

OBJECT_NAME	VARCHAR2(128)
SUBOBJECT_NAME	VARCHAR2(30)
OBJECT_ID	NUMBER
DATA_OBJECT_ID	NUMBER
OBJECT_TYPE	VARCHAR2(19)
CREATED	DATE
LAST_DDL_TIME	DATE
TIMESTAMP	VARCHAR2(19)
STATUS	VARCHAR2(7)
GENERATED	VARCHAR2(1)
SECONDARY	VARCHAR2(1)
NAMESPACE	NUMBER
EDITION_NAME	VARCHAR2(30)

```
SQL> insert into dr_object
  2  select owner, object_name, subobject_name, object_id,
  3  data_object_id, object_type, created, last_ddl_time,
  4  timestamp, status, generated, secondary, namespace, edition_name
  5  from dba_objects where object_id >= 20000;
```

52487 rows created.

```
SQL> commit;
```

Commit complete.

```
SQL> conn / as sysdba
```

Connected.

```
SQL> alter system checkpoint;
```

System altered.

导出表数据

```
[road@CentOS bin]$ ./mydul config=../etc/config.ini
```

```
MYDUL> export dict
```

```
MYDUL> load dict
```

```
MYDUL> unload table andy.dr_object
```

导入表数据

```
[road@CentOS bin]$ cd ../data
[road@CentOS data]$ ls
ANDY_DR_OBJECT.dat
[road@CentOS data]$ ../bin/myimp jake/jake file=ANDY_DR_OBJECT.dat touser=jake
```

比较数据一致性

```
SQL> conn / as sysdba
Connected.
```

```
SQL> select * from andy.dr_object minus select * from jake.dr_object;
no rows selected
```

结论

原始数据与 mydul 导出的数据一致，测试通过。

hidden 列测试

隐藏列有两种形式，一种是上面看到的 unused，查询时不显示，但是已有的列值还存储在数据段中。另一种是基于函数索引（function based index）的列，这是 Oracle 系统内部产生的列，可能需要在数据段中存储计算表达式的值，也可能不需要。

准备数据

```
SQL> conn andy/andy
Connected.
SQL> create table hi_object as select * from dba_objects where object_id<10000;

Table created.

SQL> select
  2  column_id col#, segment_column_id segcol#, internal_column_id intcol#,
  3  column_name, hidden_column
```

```
4 from dba_tab_cols where owner='ANDY' and table_name='HI_OBJECT' ;
```

COL#	SEGCOL#	INTCOL#	COLUMN_NAME	HID
1	1	1	OWNER	NO
2	2	2	OBJECT_NAME	NO
3	3	3	SUBOBJECT_NAME	NO
4	4	4	OBJECT_ID	NO
5	5	5	DATA_OBJECT_ID	NO
6	6	6	OBJECT_TYPE	NO
7	7	7	CREATED	NO
8	8	8	LAST_DDL_TIME	NO
9	9	9	TIMESTAMP	NO
10	10	10	STATUS	NO
11	11	11	TEMPORARY	NO

COL#	SEGCOL#	INTCOL#	COLUMN_NAME	HID
12	12	12	GENERATED	NO
13	13	13	SECONDARY	NO
14	14	14	NAMESPACE	NO
15	15	15	EDITION_NAME	NO

15 rows selected.

```
SQL> create index f_sub_index on hi_object (substr(object_name, 1, 8));
```

Index created.

```
SQL> select
```

```
2 column_id col#, segment_column_id segcol#, internal_column_id intcol#,
3 column_name, hidden_column
4 from dba_tab_cols where owner='ANDY' and table_name='HI_OBJECT' ;
```

COL#	SEGCOL#	INTCOL#	COLUMN_NAME	HID
1	1	1	OWNER	NO
2	2	2	OBJECT_NAME	NO
3	3	3	SUBOBJECT_NAME	NO
4	4	4	OBJECT_ID	NO
5	5	5	DATA_OBJECT_ID	NO
6	6	6	OBJECT_TYPE	NO
7	7	7	CREATED	NO

8	8	8	LAST_DDL_TIME	NO
9	9	9	TIMESTAMP	NO
10	10	10	STATUS	NO
11	11	11	TEMPORARY	NO
COL#	SEGCOL#	INTCOL#	COLUMN_NAME	HID
12	12	12	GENERATED	NO
13	13	13	SECONDARY	NO
14	14	14	NAMESPACE	NO
15	15	15	EDITION_NAME	NO
			16 SYS_NC00016\$	YES

16 rows selected.

从上面看到多出一列来，这个隐藏列就是函数索引创建的列。

导出表数据

```
[road@CentOS bin]$ ./mydul config=../etc/config.ini
```

```
MYDUL> export dict
```

```
MYDUL> load dict
```

```
MYDUL> unload table andy.hi_object
```

导入表数据

```
[road@CentOS bin]$ cd ../data
```

```
[road@CentOS data]$ ls
```

```
ANDY_HI_OBJECT.dat
```

```
[road@CentOS data]$ ../bin/myimp jake/jake file=ANDY_HI_OBJECT.dat touser=jake
```

比较数据一致性

```
SQL> conn / as sysdba
```

```
Connected.
```

```
SQL> select * from andy.hi_object minus select * from jake.hi_object;
```

```
no rows selected
```

结论

原始数据与 mydul 导出的数据一致，测试通过。

行连接测试

为了方便造成行连接，我们需要把表空间的数据块变小，设为 2K 大小。

创建表空间

由于后面还要用到这个表空间做 extent map 测试，所以把表空间的 extent 设成一致大小，并设置为最小值。

先把 db_2k_cache_size 设置打开，才能创建 2K 块大小的表空间。

```
conn / as sysdba;
alter system set db_2k_cache_size=10M;

create tablespace ts_small_block
  datafile 'ts_small_01.dbf' size 100M autoextend on
  blocksize 2048
  extent management local uniform size 10K;
```

创建用户

创建一个新用户 rock，使用上面的表空间。

```
create user rock identified by rock default tablespace ts_small_block;
```

为用户授权。

```
grant connect, resource to rock;
```

创建表

在这儿创建一个能人为产生行连接的表，让多个字段长度足够大，一个数据块不能完全存储一行数据。

```
conn rock/rock;
```

```
create table t_chain (id number, f1 varchar2(2000),  
    f2 varchar2(2000), f3 varchar2(2000));
```

插入数据

```
begin  
    for i in 1 .. 100 loop  
        insert into t_chain values (i, rpad(i, 1000, 'a'),  
            rpad(i, 1000, 'b'), rpad(i, 400, 'c'));  
    end loop;  
    commit;  
end;  
/
```

执行下面的语句，查看行连接的情况。

```
SQL> analyze table t_chain compute statistics;  
Table analyzed.
```

```
SQL> select chain_cnt from user_tables where table_name='T_CHAIN';
```

```
CHAIN_CNT  
-----  
          100
```

我们看到，所有插进去的 100 条数据都是行连接的。

导出表数据

```
[road@CentOS bin]$ ./mydul config=./etc/config.ini
```

```
MYDUL> export dict
```

```
MYDUL> load dict
```

```
MYDUL> unload table rock.t_chain
```

导入表数据

```
[road@CentOS bin]$ cd ../data
```

```
[road@CentOS data]$ ls
```

```
ROCK_T_CHAIN.dat
```

```
[road@CentOS data]$ ../bin/myimp jake/jake file=ROCK_T_CHAIN.dat touser=jake
```

比较数据一致性

```
SQL> conn / as sysdba
Connected.
SQL> select * from rock.t_chain minus select * from jake.t_chain;

no rows selected
```

结论

原始数据与 mydul 导出的数据一致，测试通过。

行迁移测试

仍然使用行连接测试中用到的 2K 数据块表空间和用户 rock 进行行迁移测试。

创建表

```
create table mig_tab (id number, f1 varchar2(2000),
  f2 varchar2(2000), f3 varchar2(2000));
```

准备数据

每行插入 800 多字节的数据，这样每块中有两行数据。

```
begin
  for i in 1 .. 100 loop
    insert into mig_tab values (i, rpad(i, 200, 'a'),
      rpad(i, 200, 'b'), rpad(i, 400, 'c'));
  end loop;
  commit;
end;
/
```

```
SQL> analyze table mig_tab compute statistics;
```


Table analyzed.

```
SQL> select chain_cnt from user_tables where table_name='MIG_TAB';
```

```
CHAIN_CNT
-----
          0
```

从上面的结果看到，没有行连接和行迁移的数据。

把一个块中的第二行数据扩大到 1200 字节，这样第二行数据就会被迁移出去。

```
begin
  for i in 1 .. 50 loop
    update mig_tab set f1=rpads(i*2, 600, 'a') where id=i*2;
  end loop;
  commit;
end;
/
```

```
SQL> analyze table mig_tab compute statistics;
```

Table analyzed.

```
SQL> select chain_cnt from user_tables where table_name='MIG_TAB';
```

```
CHAIN_CNT
-----
          50
```

在这儿看到，有 50 行数据迁移出了原来的块。

如果要看得更精确些，可以用 `analyze table mig_tab list chained rows` 语句分析后，查看 `chained_rows` 表的结果。

先创建 `chained_rows` 表

```
create table CHAINED_ROWS (
  owner_name      varchar2(30),
  table_name      varchar2(30),
  cluster_name    varchar2(30),
  partition_name  varchar2(30),
  subpartition_name varchar2(30),
  head_rowid      urowid,
  analyze_timestamp date
);
```

```
SQL> analyze table mig_tab list chained rows;
```

Table analyzed.

```
SQL> select m.id, m.rowid from mig_tab m, chained_rows c
      2   where m.rowid = c.head_rowid
      3   order by m.id;
```

ID ROWID

```
-----
      2 AAASdMAAKAAAAEAAAB
      4 AAASdMAAKAAAAAD/AAB
      6 AAASdMAAKAAAAEFAAB
      8 AAASdMAAKAAAAEBAAB
     10 AAASdMAAKAAAAECAAB
     12 AAASdMAAKAAAAEEAAB
     14 AAASdMAAKAAAAEDAAB
     16 AAASdMAAKAAAAEKAAB
     18 AAASdMAAKAAAAEGAAB
     20 AAASdMAAKAAAAEHAAB
     22 AAASdMAAKAAAAEJAAB
```

ID ROWID

```
-----
     24 AAASdMAAKAAAAEIAAB
     26 AAASdMAAKAAAAEPAAB
     28 AAASdMAAKAAAAEMAAB
     30 AAASdMAAKAAAAEOAAB
     32 AAASdMAAKAAAAENAAB
     34 AAASdMAAKAAAAEUAAB
     36 AAASdMAAKAAAAEQaab
     38 AAASdMAAKAAAAERAAB
     40 AAASdMAAKAAAAETAAB
     42 AAASdMAAKAAAAESAAB
     44 AAASdMAAKAAAAEZAAB
```

ID ROWID

```
-----
     46 AAASdMAAKAAAAEVAAB
     48 AAASdMAAKAAAAEWAAB
     50 AAASdMAAKAAAAEYAAB
     52 AAASdMAAKAAAAEXAAB
     54 AAASdMAAKAAAAEeAAB
     56 AAASdMAAKAAAAEbAAB
     58 AAASdMAAKAAAAEdAAB
```

```
60 AAASdMAAKAAAAEcAAB
62 AAASdMAAKAAAAEjAAB
64 AAASdMAAKAAAAEfAAB
66 AAASdMAAKAAAAEgAAB
```

```
ID ROWID
```

```
-----
68 AAASdMAAKAAAAEiAAB
70 AAASdMAAKAAAAEhAAB
72 AAASdMAAKAAAAEoAAB
74 AAASdMAAKAAAAEkAAB
76 AAASdMAAKAAAAElAAB
78 AAASdMAAKAAAAEnAAB
80 AAASdMAAKAAAAEmAAB
82 AAASdMAAKAAAAEtAAB
84 AAASdMAAKAAAAEqAAB
86 AAASdMAAKAAAAEsAAB
88 AAASdMAAKAAAAErAAB
```

```
ID ROWID
```

```
-----
90 AAASdMAAKAAAAEyAAB
92 AAASdMAAKAAAAEuAAB
94 AAASdMAAKAAAAEvAAB
96 AAASdMAAKAAAAExAAB
98 AAASdMAAKAAAAEwAAB
100 AAASdMAAKAAAAE3AAB
```

50 rows selected.

这样就能看到具体是哪些行发生了迁移。

别忘了 checkpoint 数据。

```
SQL> conn / as sysdba
```

```
Connected.
```

```
SQL> alter system checkpoint;
```

```
System altered.
```

导出表数据

```
[road@CentOS bin]$ ./mydul config=./etc/config.ini
```

```
MYDUL> export dict
```

```
MYDUL> load dict
```

```
MYDUL> unload table rock.mig_tab
```

导入表数据

```
[road@CentOS bin]$ cd ../data
[road@CentOS data]$ ls
ROCK_MIG_TAB.dat
[road@CentOS data]$ ../bin/myimp jake/jake file=ROCK_MIG_TAB.dat touser=jake
```

比较数据一致性

```
SQL> conn / as sysdba
Connected.
SQL> select * from rock.mig_tab minus select * from jake.mig_tab;
no rows selected
```

结论

原始数据与 mydul 导出的数据一致，测试通过。

行迁移后连接的测试

本节测试是把前面行连接和行迁移综合起来进行测试，行迁移后，再次出现行连接，行头数据会有新的变化，本节测试针对这种情况。本节测试在上一节行迁移测试的基础上进行，在上一节中，偶数行已经是行迁移状态了，我们只需把它再变成行连接就可以了。

数据准备

在上一节形成的连接行，f1 是 600 字节，f2 是 200 字节，f3 是 400 字节。如果把 f2 更新成 1000 字节，那么一行数据总共就有 2000 多字节，在加上块头的和其他空间，那么在一个块中肯定是存不下了，就会出现行连接。

从上面的描述中看到，在理论上应该会出现行迁移后再连接的情况，但在实际操作中，发现 Oracle 不希望这种情况发生，总是千方百计让数据存储在最少的块中。

试试制造一个单条的行迁移后连接的数据行。

还是用 mig_tab 表，把数据都删除掉，重新执行下面的操作，插入三行数据。

```
insert into mig_tab values (1, rpad(1, 400, 'a'), rpad(1, 400, 'b'), null);
insert into mig_tab values (2, rpad(2, 400, 'a'), rpad(2, 400, 'b'), null);
insert into mig_tab values (3, rpad(3, 400, 'a'), rpad(3, 400, 'b'), null);
commit;
```

看看这三条数据都在哪个块中。

```
select
  id,
  dbms_rowid.rowid_relative_fno(rowid) rfn,
  dbms_rowid.rowid_block_number(rowid) blk#
from mig_tab;
```

```
SQL> select
  2   id,
  3   dbms_rowid.rowid_relative_fno(rowid) rfn,
  4   dbms_rowid.rowid_block_number(rowid) blk#
  5  from mig_tab;
```

ID	RFN	BLK#
1	10	394
2	10	394
3	10	395

看到前两行数据在 394 块中，那么我们修改 id=2 的数据，使它成为迁移行。

```
update mig_tab set f2=rpad(2, 1000, 'b') where id=2;
commit;
```

把 394 块 dump 出来看一下，在这儿把主要的粘贴过来。

```
tab 0, row 1, @0x128
t1: 9 fb: --H----- lb: 0x2 cc: 0
nrid: 0x0280018e.0
end_of_block_dump
End dump data blocks tsn: 10 file#: 10 minblk 394 maxblk 394
```

看到原来的行迁移到了 0x0280018e.0 处，也就是 0x0280018e<<10>>10=398 块的第一行。

执行下面的语句，使这一行变成连接行。

```
update mig_tab set f3=rpad(2, 1000, 'c') where id=2;
commit;
```

把 398 块 dump 出来看一下。

```
tab 0, row 0, @0x1f0
tl: 1424 fb: ----F--- lb: 0x1 cc: 3
hrid: 0x0280018a.1
nrid: 0x0280018d.0
col 0: [ 2] c1 03
col 1: [400]
```

终于看到行连接了，看看行头既有 hrid 又有 nrid，正是我们希望出现的情况。先拿这三条数据测试一下吧。

导出表数据

```
[road@CentOS bin]$ ./mydul config=../etc/config.ini
```

```
MYDUL> export dict
MYDUL> load dict
MYDUL> unload table rock.mig_tab
```

导入表数据

```
[road@CentOS bin]$ cd ../data
[road@CentOS data]$ ls
ROCK_MIG_TAB.dat
[road@CentOS data]$ ../bin/myimp jake/jake file=ROCK_MIG_TAB.dat touser=jake
```

比较数据一致性

```
SQL> conn / as sysdba
Connected.
SQL> select * from rock.mig_tab minus select * from jake.mig_tab;
no rows selected
```

结论

原始数据与 mydul 导出的数据一致，测试通过。

LONG 字段表测试

本节测试还是使用块大小为 2K 的表空间，用户 rock 来进行测试，这样可以造成更多的连接块，是测试深入。测试分三种情况，第一，long 数据小于 200 字节，当做普通 varchar2 测试；第二，long 数据比较大，造成行连接，并且连接多个数据块；第三，long 字段后还有其他字段，long 数据比较大。

创建表

```
create table long_tab (id number, f1 varchar2(200), f2 varchar2(400),  
    f3 varchar2(200), f1 long);
```

插入数据

```
begin  
    for i in 1 .. 40 loop  
        insert into long_tab values (i, rpad(i, 100, 'a'), rpad(i, 100, 'b'),  
            rpad(i, 100, 'c'), rpad(i, 200, 'l'));  
    end loop;  
    commit;  
end;  
/
```

checkpoint 数据

导出表数据

```
[road@CentOS bin]$ ./mydul config=../etc/config.ini
```

```
MYDUL> export dict
```

```
MYDUL> load dict
```

```
MYDUL> unload table rock.long_tab
```

导入表数据

```
[road@CentOS bin]$ cd ../data
```

```
[road@CentOS data]$ ls
ROCK_LONG_TAB.dat
[road@CentOS data]$ ../bin/myimp jake/jake file=ROCK_LONG_TAB.dat touser=jake
```

比较数据一致性

有 long 字段的表，不能直接用 select * 查询比较了。
先来比较一下非 long 字段的数据。

```
SQL> select id, f1, f2, f3 from rock.long_tab
2 minus
3 select id, f1, f2, f3 from jake.long_tab;
```

```
no rows selected
```

用下面的 PL/SQL 检查 long 字段的一致性

```
set serveroutput on size 10000;
```

```
declare
    match boolean;
    v_id1 number;
    v_id2 number;
    v_f11 varchar2(20000);
    v_f12 varchar2(20000);
    cursor c1 is select id, f1 from rock.long_tab order by id;
    cursor c2 is select id, f1 from jake.long_tab order by id;
begin
    match := true;
    open c1;
    open c2;
    loop
        fetch c1 into v_id1, v_f11;
        exit when c1%NOTFOUND;
        fetch c2 into v_id2, v_f12;
        if v_f11 <> v_f12 then
            match := false;
            dbms_output.put_line('row id=' || v_id1 || ' long column not match. ');
        end if;
    end loop;
    close c1;
    close c2;
    if match then
        dbms_output.put_line('all long column data matched. ');
    end if;
```



```
end;
/  
  
SQL> declare  
  2  match boolean;  
  3  v_id1 number;  
  4  v_id2 number;  
  5  v_fl1 varchar2(20000);  
.....  
 27 end;  
 28 /  
all long column data matched.
```

PL/SQL procedure successfully completed.

大数据量测试

在上面的基础上，先把数据删除，然后插入数据。

```
SQL> delete from long_tab;
```

40 rows deleted.

```
SQL> commit;
```

Commit complete.

插入数据

```
begin  
  for i in 1 .. 40 loop  
    insert into long_tab values (i, rpad(i, 100, 'a'), rpad(i, 100, 'b'),  
      rpad(i, 100, 'c'), rpad(i, 30000, 'l'));  
  end loop;  
  commit;  
end;  
/
```

checkpoint 数据

用前面的方法导出数据，生成文件 ROCK_LONG_TAB.dat。

在用户 jake 下，先把 long_tab 中的数据删除

用下面的命令导入数据

```
[road@CentOS data]$ ../bin/myimp jake/jake file=ROCK_LONG_TAB.dat touser=jake
ignore=y
```

用前面的方法比较数据，先比较非 long 字段，再单独比较 long 字段数据。

```
SQL> select id, f1, f2, f3 from rock.long_tab
2 minus
3 select id, f1, f2, f3 from jake.long_tab;
```

no rows selected

```
SQL> declare
2 match boolean;
3 v_id1 number;
4 v_id2 number;
5 v_f11 varchar2(32000);
6 v_f12 varchar2(32000);
7 cursor c1 is select id, f1 from rock.long_tab order by id;
8 cursor c2 is select id, f1 from jake.long_tab order by id;
9 begin
10 match := true;
11 open c1;
12 open c2;
13 loop
14 fetch c1 into v_id1, v_f11;
15 exit when c1%NOTFOUND;
16 fetch c2 into v_id2, v_f12;
17 if v_f11 <> v_f12 then
18 match := false;
19 dbms_output.put_line('row id='||v_id1||' long column not match.');
```

```
20 end if;
21 end loop;
22 close c1;
23 close c2;
24 if match then
25 dbms_output.put_line('all long column data matched.');
```

```
26 end if;
27 end;
28 /
all long column data matched.
```

PL/SQL procedure successfully completed.

两个表中的数据一致。

long 字段不在最后

在上面测试的基础上，在表中添加两个字段，这样 LONG 字段后面就有了新字段。

准备数据：

```
alter table long_tab add (f4 varchar2(200), f5 varchar2(200));
```

插入新数据

```
begin
  for i in 41 .. 60 loop
    insert into long_tab values (i, rpad(i, 100, 'a'), rpad(i, 100, 'b'),
      rpad(i, 100, 'c'), rpad(i, 6000, 'l'), rpad(i, 40, 'd'), rpad(i, 40, 'e'));
  end loop;
  commit;
end;
/
```

checkpoint 数据

导出数据

```
[road@CentOS bin]$ ./mydul config=../etc/config.ini
```

```
MYDUL> export dict
```

```
MYDUL> load dict
```

```
MYDUL> unload table rock.long_tab
```

导入数据

```
[road@CentOS data]$ ls
```

```
ROCK_LONG_TAB.dat
```

```
[road@CentOS data]$ ../bin/myimp jake/jake file=ROCK_LONG_TAB.dat touser=jake
ignore=y
```

用前面的方法比较数据，先比较非 long 字段，再单独比较 long 字段数据。

```
SQL> select id, f1, f2, f3, f4, f5 from rock.long_tab
  2 minus
  3 select id, f1, f2, f3, f4, f5 from jake.long_tab;
```

```
no rows selected
```

比较 LONG 字段数据。

```
SQL> set serveroutput on size 10000;
```

```
SQL> declare
```

```
2  match  boolean;
3  v_id1  number;
4  v_id2  number;
5  v_f11  varchar2(32000);
6  v_f12  varchar2(32000);
7  cursor c1 is select id, f1 from rock.long_tab order by id;
8  cursor c2 is select id, f1 from jake.long_tab order by id;
9  begin
10  match := true;
11  open c1;
12  open c2;
13  loop
14  fetch c1 into v_id1, v_f11;
15  exit when c1%NOTFOUND;
16  fetch c2 into v_id2, v_f12;
17  if v_f11 <> v_f12 then
18  match := false;
19  dbms_output.put_line('row id='||v_id1||' long column not match.');
```

```
20  end if;
21  end loop;
22  close c1;
23  close c2;
24  if match then
25  dbms_output.put_line('all long column data matched.');
```

```
26  end if;
27  end;
28  /
```

```
all long column data matched.
```

PL/SQL procedure successfully completed.

结论

原始数据与 mydul 导出的数据一致，测试通过。

BASIC LOB 测试

本节测试 basicfile 的 LOB 字段表的导出。共有四个方面要测试，第一，测试 in row 的 LOB 数据；第二，chunk 小于 12 个的测试；第三，chunk 大于 12 个的测试；第四，out row LOB 的测试。

in row LOB

创建表，仍然使用 rock 用户。

```
create table in_lob (id number, f1 varchar2(200), fc clob);
```

插入数据

```
begin
  for i in 1 .. 20 loop
    insert into in_lob values (i, rpad(i, 20, 'a'), rpad(i, 600, 'l'));
  end loop;
  commit;
end;
/
```

导出数据

```
[road@CentOS bin]$ ./mydul config=../etc/config.ini
```

```
MYDUL> export dict
```

```
MYDUL> load dict
```

```
MYDUL> unload table rock.in_lob
```

导入数据

```
[road@CentOS data]$ ls
```

```
ROCK_IN_LOB.dat
```

```
[road@CentOS data]$ ../bin/myimp jake/jake file=ROCK_IN_LOB.dat touser=jake
```

比较数据

先比较非 LOB 字段

```
SQL> select id, f1 from rock.in_lob minus select id, f1 from jake.in_lob;
```

```
no rows selected
```

用下面的 PL/SQL 比较 LOB 字段

```
set serveroutput on size 10000;
```

```
declare
  match boolean;
  v_id1 number;
  v_id2 number;
  lob1 clob;
  lob2 clob;
  cursor c1 is select id, fc from rock.in_lob order by id;
```

```
cursor c2 is select id, fc from jake.in_lob order by id;
begin
  match := true;
  open c1;
  open c2;
  loop
    fetch c1 into v_id1, lob1;
    exit when c1%NOTFOUND;
    fetch c2 into v_id2, lob2;
    if dbms_lob.compare(lob1, lob2) <> 0 then
      match := false;
      dbms_output.put_line('row id=' ||v_id1||' lob column not matak.'');
    end if;
  end loop;
  close c1;
  close c2;
  if match then
    dbms_output.put_line('');
    dbms_output.put_line('all lob column data matched.'');
  end if;
end;
/
```

比较结果如下

```
SQL> declare
  2  match boolean;
  3  v_id1 number;
  4  v_id2 number;
  5  lob1 clob;
  6  lob2 clob;
  7  cursor c1 is select id, fc from rock.in_lob order by id;
  8  cursor c2 is select id, fc from jake.in_lob order by id;
  9  begin
 10  match := true;
 11  open c1;
 12  open c2;
 13  loop
 14  fetch c1 into v_id1, lob1;
 15  exit when c1%NOTFOUND;
 16  fetch c2 into v_id2, lob2;
 17  if dbms_lob.compare(lob1, lob2) <> 0 then
 18  match := false;
 19  dbms_output.put_line('row id=' ||v_id1||' lob column not matak.'');
 20  end if;
```

```
21   end loop;
22   close c1;
23   close c2;
24   if match then
25       dbms_output.put_line('');
26       dbms_output.put_line('all lob column data matched.');
```

27 end if;

```
28 end;
29 /
all lob column data matched.
```

PL/SQL procedure successfully completed.

小于 12 个 chunk

在上面创建的表中，一个 chunk 占用一个数据块，一个块是 2048 字节，减去 LOB 块的头和尾的 60 字节，每个块最多能存储 1988 个字节，12 个块最多存储 23856 字节，每个字符用两个字节表示，那么最多存储 11928 个字符，只要我们插入 LOB 的字符数少于 11928，那么 chunk 地址都在行内存储。

我们还是使用上面已经创建的 in_lob 表，把数据删除，重新插入，LOB 字符插入 8000 个。由于 LOB 字段只能直接插入 4000 个字符，所以需要用到 DBMS_LOB 包来写入 LOB 数据。

```
declare
    v_data varchar2(32767);
    v_clob clob;
begin
    for i in 1 .. 20 loop
        insert into in_lob values (i, rpad(i, 20, 'a'), empty_clob)
            returning fc into v_clob;
        v_data := rpad(i, 8000, '1');
        dbms_lob.writeappend(v_clob, 8000, v_data);
    end loop;
    commit;
end;
/
```

```
SQL> conn rock/rock
```

```
Connected.
```

```
SQL> delete from in_lob;
```

```
20 rows deleted.
```

```
SQL> commit;
```

Commit complete.

```
SQL> declare
  2   v_data varchar2(32767);
  3   v_clob clob;
  4   begin
  5     for i in 1 .. 20 loop
  6       insert into in_lob values (i, rpad(i, 20, 'a'), empty_clob)
  7         returning fc into v_clob;
  8       v_data := rpad(i, 8000, 'l');
  9       dbms_lob.writeappend(v_clob, 8000, v_data);
 10     end loop;
 11     commit;
 12 end;
 13 /
```

PL/SQL procedure successfully completed.

依照前面的步骤导出数据，然后导入到用户 jake 下

比较数据

```
SQL> select id, f1 from rock.in_lob minus select id, f1 from jake.in_lob;
```

no rows selected

```
SQL> set serveroutput on size 10000;
```

```
SQL> declare
  2   match boolean;
  3   v_id1 number;
  4   v_id2 number;
  5   lob1 clob;
  6   lob2 clob;
  7   cursor c1 is select id, fc from rock.in_lob order by id;
  8   cursor c2 is select id, fc from jake.in_lob order by id;
  9   begin
 10     match := true;
 11     open c1;
 12     open c2;
 13     loop
 14       fetch c1 into v_id1, lob1;
 15       exit when c1%NOTFOUND;
 16       fetch c2 into v_id2, lob2;
```



```
17     if dbms_lob.compare(lob1, lob2) <> 0 then
18         match := false;
19         dbms_output.put_line('row id='||v_id1||' lob column not matach.');
```

```
20     end if;
21 end loop;
22 close c1;
23 close c2;
24 if match then
25     dbms_output.put_line('');
26     dbms_output.put_line('all lob column data matched.');
```

```
27 end if;
28 end;
29 /
```

```
all lob column data matched.
```

PL/SQL procedure successfully completed.

大于 12 个 chunk

根据对小于 12 个 chunk 的 lob 存储的分析, 加大字符个数就能使 chunk 数大于 12, 需要通过 LOB index 才能完全导出 LOB 数据。

还是使用 rock 用户下的 in_lob 表, 插入 LOB 字符数为 25000 个。

```
SQL> show user
```

```
USER is "ROCK"
```

```
SQL> delete from in_lob;
```

```
20 rows deleted.
```

```
SQL> commit;
```

```
Commit complete.
```

```
SQL> declare
```

```
2     v_data varchar2(32767);
```

```
3     v_clob clob;
```

```
4     begin
```

```
5     for i in 1 .. 20 loop
```

```
6         insert into in_lob values (i, rpad(i, 20, 'a'), empty_clob)
```

```
7         returning fc into v_clob;
```

```
8         v_data := rpad(i, 25000, '1');
```

```
9      dbms_lob.writeappend(v_clob, 25000, v_data);
10     end loop;
11     commit;
12 end;
13 /
```

PL/SQL procedure successfully completed.

依照前述步骤导出和导入数据。

比较数据，结果如下：

```
SQL> select id, f1 from rock.in_lob minus select id, f1 from jake.in_lob;
```

no rows selected

```
SQL> set serveroutput on size 10000;
```

```
SQL>
```

```
SQL> declare
```

```
2     match boolean;
3     v_id1 number;
4     v_id2 number;
5     lob1 clob;
6     lob2 clob;
7     cursor c1 is select id, fc from rock.in_lob order by id;
8     cursor c2 is select id, fc from jake.in_lob order by id;
9     begin
10    match := true;
11    open c1;
12    open c2;
13    loop
14        fetch c1 into v_id1, lob1;
15        exit when c1%NOTFOUND;
16        fetch c2 into v_id2, lob2;
17        if dbms_lob.compare(lob1, lob2) <> 0 then
18            match := false;
19            dbms_output.put_line('row id='||v_id1||' lob column not matak.'');
20        end if;
21    end loop;
22    close c1;
23    close c2;
24    if match then
25        dbms_output.put_line('');
26        dbms_output.put_line('all lob column data matched.'');
27    end if;
28 end;
```

29 /

```
all lob column data matched.
```

PL/SQL procedure successfully completed.

out row LOB

仍然在 rock 用户下测试，创建 out row 的 LOB 字段表。

```
create table out_lob (id number, f1 varchar2(100), fco clob)
  lob (fco) store as (disable storage in row);
```

插入数据，在 CLOB 字段中插入 12000 个字符。

```
declare
  v_data varchar2(32767);
  v_clob clob;
begin
  for i in 1 .. 20 loop
    insert into out_lob values (i, rpad(i, 20, 'a'), empty_clob)
      returning fco into v_clob;
    v_data := rpad(i, 12000, 'l');
    dbms_lob.writeappend(v_clob, 12000, v_data);
  end loop;
  commit;
end;
/
```

依照前面的方法导出和导入数据。

比较数据一致性，结果如下：

```
SQL> select id, f1 from rock.in_lob minus select id, f1 from jake.in_lob;
```

```
no rows selected
```

```
SQL> set serveroutput on size 10000;
```

```
SQL> SQL> declare
  2  match boolean;
  3  v_id1 number;
  4  v_id2 number;
  5  lob1 clob;
  6  lob2 clob;
  7  cursor c1 is select id, fco from rock.out_lob order by id;
  8  cursor c2 is select id, fco from jake.out_lob order by id;
  9  begin
 10  match := true;
```

```
11  open c1;
12  open c2;
13  loop
14      fetch c1 into v_id1, lob1;
15      exit when c1%NOTFOUND;
16      fetch c2 into v_id2, lob2;
17      if dbms_lob.compare(lob1, lob2) <> 0 then
18          match := false;
19          dbms_output.put_line('row id='||v_id1||' lob column not matach.');
```

20 end if;

```
21  end loop;
22  close c1;
23  close c2;
24  if match then
25      dbms_output.put_line('');
26      dbms_output.put_line('all lob column data matched.');
```

27 end if;

```
28  end;
29  /
```

all lob column data matched.

PL/SQL procedure successfully completed.

Secure LOB 测试

本节测试 SecureFile LOB 字段表的导出，测试分三个部分，第一，In row LOB；第二，In row LOB 超出 4000 字节，chunk 地址在行内；第三，In row LOB 数据比较大，头块地址在行内。至于 Out row 的 LOB 存储方式与 In Row LOB 的后两种方式一样，就不测试了。

测试中继续使用 rock 用户。

in row LOB

先创建一个 SecureFile LOB 字段的表。

```
create table in_sec_lob (id number, f1 varchar2(100), fsc clob)
  lob (fsc) store as securefile;
```

插入数据

```
begin
  for i in 1 .. 20 loop
```

```
insert into in_sec_lob values (i, rpad(i, 20, 'a'), rpad(i, 400, 'l'));
end loop;
commit;
end;
/
```

这儿插入数据时产生错误，ORA-60019: Creating initial extent of size 5 in tablespace of extent size 14, 用户 rock 所在的表空间初始 extent 太小。重新建一个表空间，重建一个新用户测试 SecureFile LOB。

创建表空间 ts_secure 作为测试默认表空间

```
create tablespace ts_secure
  datafile 'ts_secure_01.dbf' size 100M autoextend on
  blocksize 2048
  extent management local uniform size 32768;
```

创建用户 secu 作为测试用户

```
create user secu identified by secu default tablespace ts_secure;
```

为 secu 用户赋权限

```
grant connect, resource to secu;
```

创建包含 SecureFile LOB 字段的测试表

```
create table in_sec_lob (id number, f1 varchar2(100), fsc clob)
  lob (fsc) store as securefile;
```

插入行内数据

```
begin
  for i in 1 .. 20 loop
    insert into in_sec_lob values (i, rpad(i, 20, 'a'), rpad(i, 400, 'l'));
  end loop;
  commit;
end;
/
```

checkpoint 数据

```
SQL> conn / as sysdba
Connected.
SQL> alter system checkpoint;
```

System altered.

导出数据，导出前要把新的数据文件 ts_secure_01.dbf 加到配置中

```
[road@CentOS bin]$ ./mydul config=../etc/config.ini
```

```
MYDUL> export dict
MYDUL> load dict
MYDUL> unload table secu.in_sec_lob
```

导入数据

```
[road@CentOS data]$ ls
SECU_IN_SEC_LOB.dat
[road@CentOS data]$ ../bin/myimp jake/jake file=SECU_IN_SEC_LOB.dat touser=jake
```

比较数据，与 Basic LOB 比较方法一样

```
SQL> conn / as sysdba
Connected.
SQL> select id, f1 from secu.in_sec_lob minus select id, f1 from jake.in_sec_lob;

no rows selected
```

比较 LOB 字段

```
SQL> set serveroutput on size 10000;

declare
    match boolean;
SQL> SQL> 2 3 v_id1 number;
4 v_id2 number;
5 lob1 clob;
6 lob2 clob;
7 cursor c1 is select id, fsc from secu.in_sec_lob order by id;
8 cursor c2 is select id, fsc from jake.in_sec_lob order by id;
9 begin
10 match := true;
11 open c1;
12 open c2;
13 loop
14 fetch c1 into v_id1, lob1;
15 exit when c1%NOTFOUND;
16 fetch c2 into v_id2, lob2;
17 if dbms_lob.compare(lob1, lob2) <> 0 then
18 match := false;
19 dbms_output.put_line('row id='||v_id1||' lob column not match. ');
20 end if;
21 end loop;
22 close c1;
23 close c2;
```

```
24  if match then
25      dbms_output.put_line('');
26      dbms_output.put_line('all lob column data matched.');
```

```
27  end if;
28  end;
29  /
all lob column data matched.
```

PL/SQL procedure successfully completed.

in row chunk

把 in_sec_lob 表中的数据删除，然后插入超过 4000 字节的数据。

准备数据

```
delete from in_sec_lob
commit;
```

```
declare
    v_data varchar2(32767);
    v_clob clob;
begin
    for i in 1 .. 20 loop
        insert into in_sec_lob values (i, rpad(i, 40, 'a'), empty_clob)
            returning fsc into v_clob;
        v_data := rpad(i, 16000, '1');
        dbms_lob.writeappend(v_clob, 16000, v_data);
    end loop;
    commit;
end;
/
```

```
SQL> conn / as sysdba
```

```
Connected.
```

```
SQL> alter system checkpoint;
```

```
System altered.
```

dump 出一个数据块来，看一下是否 LOB chunk 地址在行内。

```
select
    dbms_rowid.rowid_relative_fno(rowid) rfn,
    dbms_rowid.rowid_block_number(rowid) blk#
```

```
from secu.in_sec_lob
where id=1;
```

```
SQL> select
  2   dbms_rowid.rowid_relative_fno(rowid) rfn,
  3   dbms_rowid.rowid_block_number(rowid) blk#
  4 from secu.in_sec_lob
  5 where id=1;
```

RFN	BLK#
9	521

```
SQL> alter system dump datafile 9 block 521;
```

System altered.

```
tab 0, row 0, @0x736
tl: 98 fb: --H-FL-- lb: 0x1  cc: 3
col 0: [ 2]  c1 02
col 1: [40]
 31 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61
 61 61 61 61 61 61 61 61 61 61 61 61 61 61
col 2: [50]
 00 54 00 01 02 0c 80 80 00 02 00 00 00 01 00 00 00 0b e7 45 00 1e 40 90 00
 18 21 00 7d 00 01 02 01 02 40 02 1f 01 01 02 40 02 51 0f 01 02 40 02 61 01
LOB
```

Locator:

```
Length:      84(50)
Version:     1
Byte Length: 2
LobID: 00.00.00.01.00.00.00.0b.e7.45
Flags[ 0x02 0x0c 0x80 0x80 ]:
```

```
  Type: CLOB
  Storage: SecureFile
  Characterset Format: IMPLICIT
  Partitioned Table: No
  Options: VaringWidthReadWrite
```

SecureFile Header:

```
Length: 30
Old Flag: 0x40 [ SecureFile ]
Flag 0: 0x90 [ INODE Valid ]
```

Layers:

```
Lengths Array: INODE:24
```


INODE:

```
21 00 7d 00 01 02 01 02 40 02 1f 01 01 02 40 02 51 0f 01 02
40 02 61 01
```

从上面看到 inode 的 flag1 标志是 0x21，表示行内存储的是 chunk 地址，在后面还可以看到一共有 3 个 chunk。

导出数据

```
[road@CentOS bin]$ ./mydul config=../etc/config.ini
```

```
MYDUL> export dict
```

```
MYDUL> load dict
```

```
MYDUL> unload table secu.in_sec_lob
```

导入数据

```
[road@CentOS data]$ ../bin/myimp jake/jake file=SECU_IN_SEC_LOB.dat touser=jake
ignore=y
```

比较数据

```
SQL> conn / as sysdba
```

```
Connected.
```

```
SQL> select id, f1 from secu.in_sec_lob minus select id, f1 from jake.in_sec_lob;
```

```
no rows selected
```

```
SQL> set serveroutput on size 10000;
```

```
SQL> declare
```

```
2 match boolean;
3 v_id1 number;
4 v_id2 number;
5 lob1 clob;
6 lob2 clob;
7 cursor c1 is select id, fsc from secu.in_sec_lob order by id;
8 cursor c2 is select id, fsc from jake.in_sec_lob order by id;
9 begin
10 match := true;
11 open c1;
12 open c2;
13 loop
14 fetch c1 into v_id1, lob1;
15 exit when c1%NOTFOUND;
16 fetch c2 into v_id2, lob2;
17 if dbms_lob.compare(lob1, lob2) <> 0 then
18 match := false;
19 dbms_output.put_line('row id='||v_id1||' lob column not match.');
```

```
20     end if;
21   end loop;
22   close c1;
23   close c2;
24   if match then
25     dbms_output.put_line('');
26     dbms_output.put_line('all lob column data matched.');
```

```
27   end if;
28 end;
29 /
```

```
all lob column data matched.
```

PL/SQL procedure successfully completed.

LOB Head Block

LOB 头块的测试，我们知道当 LOB 数据大到一定程度，在 LOB locator 的行内就不再存放 chunk 地址，而是存放一个管理 LOB chunk 地址的头块地址。我们要让 LOB 数据多大才能出现头块呢？在 securefile LOB 存储时，chunk 的大小不再固定，计算有些困难，但我们可以取它的最大值，也就是我们要计算出一个数据量，当达到这个数据大小时肯定会出现 LOB 头块。下面我们来计算这个大小，我们知道 LOB chunk 中的块必须是连续的，在 oracle 中，只有 extent 中的块能保证连续性，所以一个 LOB chunk 包含的块个数最多也就是 extent 中块的个数，我们在上面创建表空间时，指定了 extent 的大小是 32768，并且每个 extent 的大小都是一样的，表空间的块大小是 2K，那么一个 extent 中块的个数=32768/2048=16 个，除去一个管理块，所以一个 LOB chunk 最多有 15 个块。由于行内存储的限制，一个行内最多能存储 6 个 chunk 地址，超过 6 个就会变成存储头块地址，那么我们能得到这个 LOB 数据的值=6*15*2048=184320 字节，由于一个块还有管理用的空间，每个 chunk 也不可能都用去 15 个块的空间，所以如果插入 184320 字节的数据，肯定就会出现 LOB 头块了，如果是 CLOB，双字节存储，那么插入 184320/2=92160 个字符可以了。我们通过下面的实验验证一下。

准备数据

```
declare
  v_data varchar2(32767);
  v_clob clob;
begin
```

```
for i in 1 .. 20 loop
  insert into in_sec_lob values (i, rpad(i, 40, 'a'), empty_clob)
  returning fsc into v_clob;
  dbms_lob.open(v_clob, DBMS_LOB.LOB_READWRITE);
  for j in 1 .. 8 loop
    v_data := rpad(i, 11520, '1');
    dbms_lob.writeappend(v_clob, 11520, v_data);
  end loop;
  dbms_lob.close(v_clob);
end loop;
commit;
end;
/
```

dump 出一行数据看一下，是否行内存储的是头块。

```
SQL> select
  2   dbms_rowid.rowid_relative_fno(rowid) rfn,
  3   dbms_rowid.rowid_block_number(rowid) blk#
  4 from secu.in_sec_lob
  5 where id=1;
```

RFN	BLK#
12	521

```
SQL> conn / as sysdba
```

Connected.

```
SQL> alter system checkpoint;
```

System altered.

```
SQL> alter system dump datafile 12 block 521;
```

System altered.

```
tab 0, row 0, @0x70b
```

```
tl: 141 fb: --H-FL-- lb: 0x1 cc: 3
```

```
col 0: [ 2] c1 02
```

```
col 1: [40]
```

```
31 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61
61 61 61 61 61 61 61 61 61 61 61 61 61 61 61
```

```
col 2: [93]
```

```
00 54 00 01 02 0c 80 80 00 02 00 00 00 01 00 00 00 0c be e5 00 49 40 90 00
43 22 00 02 d0 00 02 09 01 03 00 02 1f 01 01 03 00 02 51 0f 01 03 00 02 61
```

```
0f 01 03 00 02 35 0b 01 03 00 02 4d 03 01 03 00 02 7d 03 01 03 00 02 81 0f
01 03 00 02 71 0c 01 03 00 02 91 0f 01 03 00 02 a1 07
```

LOB

Locator:

```
Length:      84(93)
Version:     1
Byte Length: 2
LobID: 00.00.00.01.00.00.00.0c.be.e5
Flags[ 0x02 0x0c 0x80 0x80 ]:
  Type: CLOB
  Storage: SecureFile
  Characterset Format: IMPLICIT
  Partitioned Table: No
  Options: VaringWidthReadWrite
```

SecureFile Header:

```
Length: 73
Old Flag: 0x40 [ SecureFile ]
Flag 0: 0x90 [ INODE Valid ]
```

Layers:

```
Lengths Array: INODE:67
```

INODE:

```
22 00 02 d0 00 02 09 01 03 00 02 1f 01 01 03 00 02 51 0f 01
03 00 02 61 0f 01 03 00 02 35 0b 01 03 00 02 4d 03 01 03 00
02 7d 03 01 03 00 02 81 0f 01 03 00 02 71 0c 01 03 00 02 91
0f 01 03 00 02 a1 07
```

从上面看到，行内存储的还是 chunk 地址，不是头块地址，为什么呢？原来我们这个表创建时，LOB 是 Enable store in row 的，这样这个字段一共能存储 4000 字节，去掉 Lob locator, Secure File Header 和 Inode 还能存储三千多字节，Oracle 在 SecureFile LOB 存储时改变了策略，剩下的这三千多字节都可以用来存储 chunk 地址，这样在一定的数据量范围内就不用通过头块索引了，减小了 I/O，真是个聪明的做法。

如果用 in row 的 LOB，为了测试 LOB 头块读取数据，就要插入大量的 LOB 数据才能出现，为了尽量用少的数据创建测试场景，我们也只有改变策略了，用 out row 的 LOB，那么我们上面分析的就生效了。

创建一个 out row 的 LOB 表

```
create table out_sec_lob (id number, f1 varchar2(100), fsc clob)
  lob (fsc) store as securefile (disable storage in row);
```

插入数据

```
declare
```

```
v_data varchar2(32767);
v_clob clob;
begin
  for i in 1 .. 20 loop
    insert into out_sec_lob values (i, rpad(i, 40, 'a'), empty_clob)
      returning fsc into v_clob;
    dbms_lob.open(v_clob, DBMS_LOB.LOB_READWRITE);
    for j in 1 .. 8 loop
      v_data := rpad(i, 11520, '1');
      dbms_lob.writeappend(v_clob, 11520, v_data);
    end loop;
    dbms_lob.close(v_clob);
  end loop;
  commit;
end;
/
```

dump 出来检查一下。

```
select
  dbms_rowid.rowid_relative_fno(rowid) rfn,
  dbms_rowid.rowid_block_number(rowid) blk#
from secu.out_sec_lob
where id=1;
```

```
SQL> conn / as sysdba
Connected.
SQL> alter system checkpoint;
```

System altered.

```
SQL> select
  2  dbms_rowid.rowid_relative_fno(rowid) rfn,
  3  dbms_rowid.rowid_block_number(rowid) blk#
  4  from secu.out_sec_lob
  5  where id=1;
```

RFN	BLK#
12	2905

```
SQL> alter system dump datafile 12 block 2905;
```

System altered.

```
tab 0, row 0, @0x744
tl: 84 fb: --H-FL-- lb: 0x1 cc: 3
col 0: [ 2] c1 02
col 1: [40]
 31 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61 61
 61 61 61 61 61 61 61 61 61 61 61 61 61 61
col 2: [36]
 00 54 00 01 02 0c 80 80 00 02 00 00 00 01 00 00 00 0c bf 17 00 10 40 90 00
 0a 42 00 02 d0 00 02 03 00 0c 21
LOB
Locator:
  Length:      84(36)
  Version:     1
  Byte Length: 2
  LobID: 00.00.00.01.00.00.00.0c.bf.17
  Flags[ 0x02 0x0c 0x80 0x80 ]:
    Type: CLOB
    Storage: SecureFile
    Characterset Format: IMPLICIT
    Partitioned Table: No
    Options: VaringWidthReadWrite
SecureFile Header:
  Length: 16
  Old Flag: 0x40 [ SecureFile ]
  Flag 0: 0x90 [ INODE Valid ]
  Layers:
    Lengths Array: INODE:10
    INODE:
      42 00 02 d0 00 02 03 00 0c 21
```

这次我们看到行内存储的是头块了，地址是 0x03000c21。

导出数据

```
[road@CentOS bin]$ ./mydul config=../etc/config.ini
MYDUL> export dict
MYDUL> load dict
MYDUL> unload table secu.out_sec_lob
```

导入数据

```
[road@CentOS bin]$ cd ../data
[road@CentOS data]$ ls
SECU_OUT_SEC_LOB.dat
```

```
[road@CentOS data]$ ../bin/myimp jake/jake file=SECU_OUT_SEC_LOB.dat touser=jake
```

比较数据一致性

先比较非 LOB 数据

```
SQL> select id, f1 from secu.out_sec_lob minus select id, f1 from jake.out_sec_lob;
```

```
no rows selected
```

再比较 LOB 数据

```
set serveroutput on size 10000;
```

```
declare
  match boolean;
  v_id1 number;
  v_id2 number;
  lob1 clob;
  lob2 clob;
  cursor c1 is select id, fsc from secu.out_sec_lob order by id;
  cursor c2 is select id, fsc from jake.out_sec_lob order by id;
begin
  match := true;
  open c1;
  open c2;
  loop
    fetch c1 into v_id1, lob1;
    exit when c1%NOTFOUND;
    fetch c2 into v_id2, lob2;
    if dbms_lob.compare(lob1, lob2) <> 0 then
      match := false;
      dbms_output.put_line('row id=' || v_id1 || ' lob column not matach. ');
    end if;
  end loop;
  close c1;
  close c2;
  if match then
    dbms_output.put_line(' ');
    dbms_output.put_line('all lob column data matched. ');
  end if;
end;
/
```

比较结果

```
SQL> declare
  2  match boolean;
```

```
3  v_id1  number;
4  v_id2  number;
5  lob1   clob;
6  lob2   clob;
7  cursor c1 is select id, fsc from secu.out_sec_lob order by id;
8  cursor c2 is select id, fsc from jake.out_sec_lob order by id;
9  begin
10 match := true;
11 open c1;
12 open c2;
13 loop
14     fetch c1 into v_id1, lob1;
15     exit when c1%NOTFOUND;
16     fetch c2 into v_id2, lob2;
17     if dbms_lob.compare(lob1, lob2) <> 0 then
18         match := false;
19         dbms_output.put_line('row id='||v_id1||' lob column not match.');
```

```
20     end if;
21 end loop;
22 close c1;
23 close c2;
24 if match then
25     dbms_output.put_line('');
26     dbms_output.put_line('all lob column data matched.');
```

PL/SQL procedure successfully completed.

分区表测试

本节测试分区表的导出，分区表有几种，我们只测试最常用的范围（range）分区表，其他种类分区表的存储与范围分区表一样，只是分区的方法不同而已，不影响导出测试。

本节测试中使用用户 andy。

准备数据

创建分区表


```
create table range_part
  (id number, f1 varchar2(1000), f2 varchar2(1000), f3 varchar2(1000))
partition by range (id)
(
  partition r_part1 values less than (1000),
  partition r_part2 values less than (2000),
  partition r_part3 values less than (3000),
  partition r_part4 values less than (maxvalue)
);
```

插入数据

```
begin
  for i in 1 .. 200 loop
    insert into range_part values
      (i, rpad(i, 100, 'a'), rpad(i, 200, 'b'), rpad(i, 300, 'c'));
  end loop;
  commit;

  for i in 1001 .. 1200 loop
    insert into range_part values
      (i, rpad(i, 100, 'a'), rpad(i, 200, 'b'), rpad(i, 300, 'c'));
  end loop;
  commit;

  for i in 2001 .. 2200 loop
    insert into range_part values
      (i, rpad(i, 100, 'a'), rpad(i, 200, 'b'), rpad(i, 300, 'c'));
  end loop;
  commit;

  for i in 3001 .. 3200 loop
    insert into range_part values
      (i, rpad(i, 100, 'a'), rpad(i, 200, 'b'), rpad(i, 300, 'c'));
  end loop;
  commit;
end;
/
```

```
SQL> conn / as sysdba
```

```
Connected.
```

```
SQL> alter system checkpoint;
```

```
System altered.
```

导出数据

```
[road@CentOS bin]$ ./mydul config=../etc/config.ini
MYDUL> export dict
MYDUL> load dict
MYDUL> unload table andy.range_part
```

导入数据

```
[road@CentOS bin]$ cd ../data
[road@CentOS data]$ ls
ANDY_RANGE_PART.dat
[road@CentOS data]$ ../bin/myimp jake/jake file=ANDY_RANGE_PART.dat touser=jake
```

比较数据

```
SQL> select * from andy.range_part minus select * from jake.range_part;

no rows selected
```

子分区表测试

准备数据

创建子分区表

```
create table comp_part
  (lid number, rid number, f1 varchar2(1000), f2 varchar2(1000), f3 varchar2(1000))
partition by list (lid)
subpartition by range (rid)
(
  partition l_part1 values (100)
  (
    subpartition r_subpart11 values less than (1000),
    subpartition r_subpart12 values less than (2000),
    subpartition r_subpart13 values less than (maxvalue)
  ),
  partition l_part2 values (200)
  (
```

```
    subpartition r_subpart21 values less than (3000),
    subpartition r_subpart22 values less than (4000),
    subpartition r_subpart23 values less than (maxvalue)
  )
);
```

插入数据

```
begin
  --part 1, subpart 1
  for i in 1 .. 200 loop
    insert into comp_part values
      (100, i, rpad(i, 100, 'a'), rpad(i, 200, 'b'), rpad(i, 300, 'c'));
  end loop;
  commit;

  -- part 1, subpart 2
  for i in 1001 .. 1200 loop
    insert into comp_part values
      (100, i, rpad(i, 100, 'a'), rpad(i, 200, 'b'), rpad(i, 300, 'c'));
  end loop;
  commit;

  -- part 1, subpart 3
  for i in 2001 .. 2200 loop
    insert into comp_part values
      (100, i, rpad(i, 100, 'a'), rpad(i, 200, 'b'), rpad(i, 300, 'c'));
  end loop;
  commit;

  --part 2, subpart 1
  for i in 1 .. 200 loop
    insert into comp_part values
      (200, i, rpad(i, 100, 'a'), rpad(i, 200, 'b'), rpad(i, 300, 'c'));
  end loop;
  commit;

  -- part 2, subpart 2
  for i in 3001 .. 3200 loop
    insert into comp_part values
      (200, i, rpad(i, 100, 'a'), rpad(i, 200, 'b'), rpad(i, 300, 'c'));
  end loop;
  commit;

  -- part 2, subpart 3
```

```
for i in 4001 .. 4200 loop
  insert into comp_part values
    (200, i, rpad(i, 100, 'a'), rpad(i, 200, 'b'), rpad(i, 300, 'c'));
end loop;
commit;
end;
/
```

```
SQL> conn / as sysdba
Connected.
SQL> alter system checkpoint;
```

System altered.

导出数据

```
[road@CentOS bin]$ ./mydul config=../etc/config.ini
MYDUL> export dict
MYDUL> load dict
MYDUL> unload table andy.comp_part
```

导入数据

```
[road@CentOS data]$ ls
ANDY_COMP_PART.dat
[road@CentOS data]$ ../bin/myimp jake/jake file=ANDY_COMP_PART.dat touser=jake
```

比较数据

```
SQL> select * from andy.comp_part minus select * from jake.comp_part;

no rows selected
```

国际字符集测试

准备数据

```
create table tab_in
  (id number, f1 char(40), f2 nchar(40), f3 varchar2(100), f4 nvarchar2(100));

insert into tab_in values
  (1, '中文测试 1', '中文测试 2', '中文测试 3', '中文测试 4');
commit;
```

导出数据

```
[road@CentOS bin]$ ./mydul config=../etc/config.ini
MYDUL> export dict
MYDUL> load dict
MYDUL> unload table andy.tab_in
```

导入数据

```
[road@CentOS data]$ ls
ANDY_TAB_IN.dat
[road@CentOS data]$ ../bin/myimp jake/jake file=ANDY_TAB_IN.dat touser=jake
ignore=y
```

比较数据

```
select id, trim(f1), trim(f2), f3, f4 from andy.tab_in
minus
select id, trim(f1), trim(f2), f3, f4 from jake.tab_in;

SQL> select id, trim(f1), trim(f2), f3, f4 from andy.tab_in
  2  minus
  3  select id, trim(f1), trim(f2), f3, f4 from jake.tab_in;

no rows selected
```

```
SQL> select id, trim(f1), trim(f2), f3, f4 from jake.tab_in;
```

```
          ID TRIM(F1)
-----
TRIM(F2)
-----
F3
-----
F4
-----
          1 中文测试 1
中文测试 2
中文测试 3
中文测试 4
```

测试 CLOB 和 NCLOB 的中文

```
create table tab_in_lob (id number, f1 clob, f2 nclob);
```

```
insert into tab_in_lob values (1, '这是一条中文字符的 LOB 数据', '这是一条 national
LOB 的中文数据');
```

```
commit;
```

导出数据

```
[road@CentOS bin]$ ./mydul config=../etc/config.ini
MYDUL> export dict
MYDUL> load dict
MYDUL> unload table andy.tab_in_lob
```

导入数据

```
[road@CentOS data]$ ls
ANDY_TAB_IN_LOB.dat
[road@CentOS data]$ ../bin/myimp jake/jake file=ANDY_TAB_IN_LOB.dat touser=jake
ignore=y
```

比较数据

不能用 minus 直接比较数据，我们查询出来看看，人工比对一下吧。

```
SQL> select * from andy.tab_in_lob;
```

```
          ID
-----
F1
-----
```

F2

1

这是一条中文字符的 LOB 数据

这是一条 national LOB 的中文数据

```
SQL> select * from jake.tab_in_lob;
```

ID

F1

F2

1

这是一条中文字符的 LOB 数据

这是一条 national LOB 的中文数据

可见，数据是一致的。

ASSM Extent Map 测试

本节测试是为了在表的数据量比较大时，Segment Header 块的 extent map 用完后，要使用下一个 extent map 来存储 extent 的地址。本节测试要验证 extent map 的读取没有问题。

我们使用 ts_small_block 表空间，rock 用户来做测试，这个表空间使用 2K 的块大小，一个 extent 包含 5 个块，所以不用很大的数据量，就能出现下一个 extent map。

准备数据

创建表

```
create table tab_extmap (id number, f1 varchar2(2000));
```

插入数据

我们在每行数据的 f1 字段插入 1600 字符，这样就让每一行数据占用一个块，我们来粗略计算一下需要插入多少行数据，才能出现下一个 extent map。一个块 2048 字节，一半的空间用于辅助 map，每个 extent rdba 占用 4 字节，那么一个块最多能存储的 extent 个数是 $2048/2/4=256$ ，一个 extent 包含 5 个块，那么一共能管理 $256*5=1280$ 个块，也就是说

插入 1280 行数据就能出现 extent map, 我们来试一下。

```
begin
  for i in 1 .. 1280 loop
    insert into tab_extmap values (i, rpad(i, 1600, 'a'));
  end loop;
  commit;
end;
/
```

```
SQL> conn / as sysdba
Connected.
SQL> alter system checkpoint;
```

System altered.

我们把段头的块 dump 出来看看。

```
SQL> select header_file, header_block from dba_segments where owner='ROCK' and
segment_name='TAB_EXTMAP' ;
```

```
HEADER_FILE HEADER_BLOCK
-----
                8          34
```

```
SQL> alter system dump datafile 8 block 34;
```

System altered.

Block dump from disk:

buffer tsn: 9 rdba: 0x02000022 (8/34)

scn: 0x0000.0040de88 seq: 0x03 flg: 0x04 tail: 0xde882303

frmt: 0x02 chkval: 0x709a type: 0x23=PAGETABLE SEGMENT HEADER

Hex dump of block: st=0, typ_found=1

Extent Control Header

```
-----
Extent Header:: spare1: 0      spare2: 0      #extents: 267      #blocks: 1335
```

```
                last map 0x0200038b #maps: 2      offset: 668
```

```
                Highwater:: 0x02000557 ext#: 266      blk#: 5      ext size: 5
```

```
#blocks in seg. hdr's freelists: 0
```

```
#blocks below: 1283
```

```
mapblk 0x0200038b offset: 91
```

```
                Unlocked
```

```
-----
Low HighWater Mark :
```



```
Highwater:: 0x02000557 ext#: 266 blk#: 5 ext size: 5
#blocks in seg. hdr's freelists: 0
#blocks below: 1283
mapblk 0x0200038b offset: 91
Level 1 BMB for High HWM block: 0x0200052a
Level 1 BMB for Low HWM block: 0x0200052a
-----
Segment Type: 1 nl2: 1 blksz: 2048 fbsz: 0
L2 Array start offset: 0x00000434
First Level 3 BMB: 0x00000000
L2 Hint for inserts: 0x02000021
Last Level 1 BMB: 0x0200052a
Last Level II BMB: 0x02000021
Last Level III BMB: 0x00000000
Map Header:: next 0x02000120 #extents: 51 obj#: 76009 flag: 0x10000000
Inc # 0
```

从上面看到下一个 extent map 的地址是 0x02000120，转换后就是相对文件号 8，块号 288 的地址，再 dump 出来看一下。

```
Block dump from disk:
buffer tsn: 9 rdba: 0x02000120 (8/288)
scn: 0x0000.0040ddce seq: 0x01 flg: 0x04 tail: 0xddce2401
frmt: 0x02 chkval: 0x7afd type: 0x24=PAGEABLE EXTENT MAP BLOCK
Hex dump of block: st=0, typ_found=1
EMB Dump:
Map Header:: next 0x0200038b #extents: 124 obj#: 76009 flag: 0x10000000
Inc # 0
```

这里看到还有下一个 extent map，原来我们在计算时犯了一个小错误，一个 extent map 不止有 rdba，还有一个 4 字节的长度，那么一个块中能存储的最大 extent 个数就是 $2048/2/8=128$ 个，实际我们看到是 #extents: 124，相去不远。

那么我们就拿这批数据来测试吧。

导出数据

```
[road@CentOS bin]$ ./mydul config=../etc/config.ini
MYDUL> export dict
MYDUL> load dict
MYDUL> unload table rock.tab_extmap
```

导入数据

```
[road@CentOS data]$ ls
ANDY_TAB_IN_LOB.dat
[road@CentOS data]$ ls
ROCK_TAB_EXTMAP.dat
[road@CentOS data]$ ../bin/myimp jake/jake file=ROCK_TAB_EXTMAP.dat touser=jake
```

比较数据

```
SQL> select * from rock.tab_extmap minus select * from jake.tab_extmap;

no rows selected
```

MSSM Extent Map 测试

MSSM 的段头块，我们在导出数据字典的时候就已经遇到了，为了向下兼容，Oracle 数据字典的表空间还是手动管理的。我们在这里主要是测试 extent map 的导出。

准备环境

创建表空间

```
create tablespace ts_mssm
  datafile 'ts_mssm_01.dbf' size 100M autoextend on
  blocksize 2048
  extent management local uniform size 10K
  segment space management manual;
```

创建用户

```
create user uman identified by uman default tablespace ts_mssm;
```

为用户授权

```
grant connect, resource to uman;
```

准备数据

创建测试表

```
create table man_extmap (id number, f1 varchar2(2000));
```

插入数据

```
begin
  for i in 1 .. 1280 loop
    insert into man_extmap values (i, rpad(i, 1600, 'a'));
  end loop;
  commit;
end;
/
```

```
SQL> conn / as sysdba;
Connected.
SQL> alter system checkpoint;
```

System altered.

检查是否出现了下一个 extent map

```
SQL> select header_file, header_block from dba_segments
  2   where owner='UMAN' and segment_name='MAN_EXTMAP';
```

```
HEADER_FILE HEADER_BLOCK
```

```
-----
          10          32
```

```
SQL> alter system dump datafile 10 block 32;
```

System altered.

看一下 dump 出来的内容。

Block dump from disk:

buffer tsn: 11 rdba: 0x02800020 (10/32)

scn: 0x0000.0040e813 seq: 0x03 flg: 0x04 tail: 0xe8131003

frmt: 0x02 chkval: 0x1f74 type: 0x10=DATA SEGMENT HEADER - UNLIMITED

Hex dump of block: st=0, typ_found=1

Extent Control Header

```
-----
Extent Header:: spare1: 0      spare2: 0      #extents: 257  #blocks: 1283
```

```
                last map 0x0280027d #maps: 1      offset: 1056
```

```
                Highwater:: 0x02800525 ext#: 256   blk#: 5      ext size: 5
```

```
#blocks in seg. hdr's freelists: 4
```

```
#blocks below: 1283
```

```
mapblk 0x0280027d offset: 135
```

```
                Unlocked
```

```
                Map Header:: next 0x0280027d #extents: 121  obj#: 76011  flag: 0x40000000
```

从上面看到，已经出现了下一个 extent map，用这批数据测试。

导出数据

```
[road@CentOS bin]$ ./mydul config=../etc/config.ini
MYDUL> export dict
MYDUL> load dict
MYDUL> unload table uman.man_extmap
```

导入数据

```
[road@CentOS data]$ ls
UMAN_MAN_EXTMAP.dat
[road@CentOS data]$ ../bin/myimp jake/jake file=UMAN_MAN_EXTMAP.dat touser=jake
```

比较数据

```
SQL> select * from uman.man_extmap minus select * from jake.man_extmap;

no rows selected
```